

# Exam questions

## Reasoning Lectures

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### 1 Questions

1. Show that the De Morgan laws  $A \wedge B \dashv\vdash \neg(\neg A \vee \neg B)$  are derivable in Classical Logic, e.g exhibit a derivation of the sequent  $A \wedge B \vdash \neg(\neg A \vee \neg B)$  and a derivation of the sequent  $\neg(\neg A \vee \neg B) \vdash A \wedge B$  in G.CL.
2. Only one of the previous sequents is derivable in Intuitionistic Logic: exhibit a derivation of it in G.IL.
3. Show that the sequent  $A \rightarrow (B \wedge C) \vdash (A \rightarrow B) \wedge (A \rightarrow C)$  is derivable in Intuitionistic Logic using G.IL.
4. The sequent  $A \multimap (B \otimes C) \vdash (A \multimap B) \otimes (A \multimap C)$  is not derivable in Intuitionistic Linear Logic. Find a variant of the previous sequent where one or more (sub)formulas appear under the scope of  $!$ , and such that it is derivable in Intuitionistic Linear Logic. Show a derivation of it in G.ILL.

### 2 Answers

1.

$$\frac{\frac{\frac{A \vdash A}{A, B \vdash A} \text{wL} \quad \frac{\frac{B \vdash B}{B, A \vdash B} \text{wL}}{A, B \vdash B} \text{eL}}{\frac{A \wedge B \vdash A}{A \wedge B, \neg A \vdash} \wedge_L \quad \frac{A \wedge B \vdash B}{A \wedge B, \neg B \vdash} \wedge_L} \neg_L \quad \frac{A \wedge B, \neg A \vdash \quad A \wedge B, \neg B \vdash}{A \wedge B, \neg A \vee \neg B \vdash} \vee_L}{A \wedge B \vdash \neg(\neg A \vee \neg B)} \neg_R$$

$$\begin{array}{c}
\frac{A \vdash A}{A, B \vdash A} \text{w}_L \quad \frac{B \vdash B}{B, A \vdash B} \text{w}_L}{\frac{A, B \vdash A}{A, B \vdash A \wedge B} \wedge_R} \text{e}_L \\
\frac{A \vdash \neg B, A \wedge B}{\vdash \neg A, \neg B, A \wedge B} \neg_R}{\frac{\vdash \neg A \vee \neg B, A \wedge B}{\neg(\neg A \vee \neg B) \vdash A \wedge B} \neg_L} \vee_R
\end{array}$$

2. The first proof in 1. is a proof in G.II as well:

$$\begin{array}{c}
\frac{A \vdash A}{A, B \vdash A} \text{w}_L \quad \frac{B \vdash B}{B, A \vdash B} \text{w}_L}{\frac{A \wedge B \vdash A}{A \wedge B, \neg A \vdash} \neg_L} \wedge_L \quad \frac{A \wedge B \vdash B}{A \wedge B, \neg B \vdash} \neg_L} \wedge_L \\
\frac{A \wedge B, \neg A \vee \neg B \vdash}{A \wedge B \vdash \neg(\neg A \vee \neg B)} \neg_R} \vee_L
\end{array}$$

3.

$$\begin{array}{c}
\frac{\frac{A \vdash A}{A, A \rightarrow (B \wedge C) \vdash B} \rightarrow_L \quad \frac{\frac{B \vdash B}{B, C \vdash B} \text{w}_L}{B \wedge C \vdash B} \wedge_L}{A \rightarrow (B \wedge C) \vdash A \rightarrow B} \wedge_L \quad \frac{\frac{A \vdash A}{A, A \rightarrow (B \wedge C) \vdash C} \rightarrow_L \quad \frac{\frac{C \vdash C}{C, B \vdash C} \text{w}_L}{B, C \vdash C} \text{e}_L}{B \wedge C \vdash C} \wedge_L}{A \rightarrow (B \wedge C) \vdash (A \rightarrow B) \wedge (A \rightarrow C)} \wedge_R} \rightarrow_L
\end{array}$$

4.

$$\begin{array}{c}
\frac{\frac{A \vdash A}{A, A \multimap (!B \otimes !C) \vdash B} \multimap_L \quad \frac{\frac{B \vdash B}{!B \vdash B} \text{d!}}{!B, !C \vdash B} \text{!w}}{A \multimap (!B \otimes !C) \vdash A \multimap B} \otimes_L}{\frac{A \multimap (!B \otimes !C) \vdash A \multimap B}{!(A \multimap (!B \otimes !C)) \vdash A \multimap B} \text{d!}} \otimes_L \quad \frac{\frac{A \vdash A}{A, A \multimap (!B \otimes !C) \vdash C} \multimap_L \quad \frac{\frac{C \vdash C}{!C \vdash C} \text{d!}}{!C, !B \vdash C} \text{!w}}{!B, !C \vdash C} \text{e}}{A \multimap (!B \otimes !C) \vdash A \multimap C} \otimes_L}{\frac{!(A \multimap (!B \otimes !C)), !(A \multimap (!B \otimes !C)) \vdash (A \multimap B) \otimes (A \multimap C)}{!(A \multimap (!B \otimes !C)) \vdash (A \multimap B) \otimes (A \multimap C)} \otimes_R} \text{!c}
\end{array}$$

### 3 Gentzen sequent calculi

#### 3.1 Sequent calculus G.CL (Classical Logic)

*Axioms*

$$\frac{}{p \vdash p} \text{Id} \quad \perp \quad \frac{}{\perp \vdash} \quad \frac{}{\vdash \top} \top$$

*Logical Rules*

$$\begin{array}{c} \wedge_L \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \quad \frac{\Gamma \vdash A, \Delta \quad \Gamma \vdash B, \Delta}{\Gamma \vdash A \wedge B, \Delta} \wedge_R \quad \neg_L \frac{\Gamma \vdash A, \Delta}{\Gamma, \neg A \vdash \Delta} \\ \\ \frac{\Gamma \vdash A, B, \Delta}{\Gamma \vdash A \vee B, \Delta} \vee_R \quad \vee_L \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \vee B \vdash \Delta} \quad \neg_R \frac{\Gamma, A \vdash \Delta}{\Gamma \vdash \neg A, \Delta} \\ \\ \rightarrow_L \frac{\Gamma \vdash A, \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \rightarrow B \vdash \Delta} \quad \rightarrow_R \frac{\Gamma, A \vdash B, \Delta}{\Gamma \vdash A \rightarrow B, \Delta} \end{array}$$

*Structural Rules*

$$\begin{array}{c} w_L \frac{\Gamma \vdash \Delta}{\Gamma, A \vdash \Delta} \quad \frac{\Gamma \vdash \Delta}{\Gamma \vdash A, \Delta} w_R \quad c_L \frac{\Gamma, A, A \vdash \Delta}{\Gamma, A \vdash \Delta} \quad \frac{\Gamma \vdash A, A, \Delta}{\Gamma \vdash A, \Delta} c_R \\ \\ e_L \frac{\Gamma, B, A, \Gamma' \vdash \Delta}{\Gamma, A, B, \Gamma' \vdash \Delta} \quad \frac{\Gamma \vdash \Delta, B, A, \Delta'}{\Gamma \vdash \Delta, A, B, \Delta'} e_R \quad \frac{\Gamma \vdash A, \Delta \quad \Gamma, A \vdash \Delta}{\Gamma \vdash \Delta} \text{cut} \end{array}$$

#### 3.2 Sequent calculus G.IL (Intuitionistic Logic)

*Axioms*

$$\frac{}{p \vdash p} \text{Id} \quad \perp \quad \frac{}{\perp \vdash} \quad \frac{}{\vdash \top} \top$$

*Logical Rules*

$$\begin{array}{c} \wedge_L \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \quad \frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \wedge B} \wedge_R \\ \\ \frac{\Gamma \vdash A}{\Gamma \vdash A \vee B} \vee_{R1} \quad \frac{\Gamma \vdash B}{\Gamma \vdash A \vee B} \vee_{R2} \quad \vee_L \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \vee B \vdash \Delta} \\ \\ \rightarrow_L \frac{\Gamma \vdash A \quad \Gamma, B \vdash \Delta}{\Gamma, A \rightarrow B \vdash \Delta} \quad \rightarrow_R \frac{\Gamma, A \vdash B}{\Gamma \vdash A \rightarrow B} \quad \neg_L \frac{\Gamma \vdash A}{\Gamma, \neg A \vdash} \quad \frac{\Gamma, A \vdash}{\Gamma \vdash \neg A} \neg_R \end{array}$$

*Structural Rules*

$$\begin{array}{c}
w_L \frac{\Gamma \vdash \Delta}{\Gamma, A \vdash \Delta} \quad \frac{\Gamma \vdash \Delta}{\Gamma \vdash A} w_R \quad c_L \frac{\Gamma, A, A \vdash \Delta}{\Gamma, A \vdash \Delta} \\
e_L \frac{\Gamma, B, A, \Gamma' \vdash \Delta}{\Gamma, A, B, \Gamma' \vdash \Delta} \quad \frac{\Gamma \vdash A \quad \Gamma, A \vdash \Delta}{\Gamma \vdash \Delta} \text{cut}
\end{array}$$

### 3.3 Sequent calculus G.ILL (Intuitionistic Linear Logic)

*Axioms + 1<sub>L</sub>*

$$\frac{}{p \vdash p} \text{Id} \quad \perp_L \frac{}{\Gamma, \perp \vdash A} \quad \frac{}{\Gamma \vdash \top} \top_R \quad \frac{}{\vdash 1} 1_R \quad 1_L \frac{\Gamma \vdash \Delta}{\Gamma, 1 \vdash \Delta}$$

*Logical Rules*

$$\begin{array}{c}
\wedge_{L1} \frac{\Gamma, A \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \quad \wedge_{L2} \frac{\Gamma, B \vdash \Delta}{\Gamma, A \wedge B \vdash \Delta} \quad \frac{\Gamma \vdash A \quad \Gamma \vdash B}{\Gamma \vdash A \wedge B} \wedge_R \\
\vee_{R1} \frac{\Gamma \vdash A}{\Gamma \vdash A \vee B} \vee_{R2} \frac{\Gamma \vdash B}{\Gamma \vdash A \vee B} \quad \vee_L \frac{\Gamma, A \vdash \Delta \quad \Gamma, B \vdash \Delta}{\Gamma, A \vee B \vdash \Delta} \\
\otimes_L \frac{\Gamma, A, B \vdash \Delta}{\Gamma, A \otimes B \vdash \Delta} \quad \frac{\Gamma \vdash A \quad \Gamma' \vdash B}{\Gamma, \Gamma' \vdash A \otimes B} \otimes_R \quad d!_L \frac{\Gamma, A \vdash \Delta}{\Gamma, !A \vdash \Delta} \quad \frac{! \Gamma \vdash A}{! \Gamma \vdash !A} p!_R \\
\multimap_L \frac{\Gamma \vdash A \quad \Gamma', B \vdash \Delta}{\Gamma, \Gamma', A \multimap B \vdash \Delta} \quad \frac{\Gamma, A \vdash B}{\Gamma \vdash A \multimap B} \multimap_R \quad \sim_L \frac{\Gamma \vdash A}{\Gamma, \sim A \vdash} \quad \frac{\Gamma, A \vdash}{\Gamma \vdash \sim A} \sim_R
\end{array}$$

*Structural Rules*

$$!w \frac{\Gamma \vdash \Delta}{\Gamma, !A \vdash \Delta} \quad !c \frac{\Gamma, !A, !A \vdash \Delta}{\Gamma, !A \vdash \Delta} \quad e \frac{\Gamma, B, A, \Gamma' \vdash \Delta}{\Gamma, A, B, \Gamma' \vdash \Delta} \quad \frac{\Gamma \vdash A \quad \Gamma', A \vdash \Delta}{\Gamma, \Gamma' \vdash \Delta} \text{cut}$$